Experiments in Why3

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Plan

- Why3
- demos

Goal

conclusions

Write elegant programs

with elegant correctness proofs

+ training in program proofs

Why3





Why3 (1/8)

A programming language tells you what a program does, Why3 tells you why it works.

- 3rd release of system Why
- developed at LRI (orsay) + Inria
- http://why3.lri.fr

```
[Jean-Christophe Filliâtre,
Claude Marché,
Andrei Paskevich,
Guillaume Melquiond,
Vincent Bolot,
et al]
```

Why3 (2/8)

- small Pascal-like imperative programming language
 - [with ML syntax 🢾 !!]
- invariants + assertions in Hoare logic

[+ recursive functions, inductive datatypes, inductive predicates]

interfaces with modern SMT's

[alt-ergo, cvc3, cvc4, eprover, gappa, simplify, spass, yices, z3]

interfaces with interactive proof assistants

[coq, pvs, isabelle-hol?]

Why3 (3/8)

a

programming language MLW

```
let swap (a: array int) (i: int) (j: int) =
 let v = a[i] in
  a[i] <- a[j];</pre>
  a[j] <- v
let selection_sort (a: array int) =
   for i = 0 to length a - 1 do
     let imin = ref i in
     for j = i + 1 to length a - 1 do
       if a[j] < a[!imin] then imin := j</pre>
     done;
     swap a !imin i
   done
                        imin
0
```

Why3 (4/8)

• Hoare logic

```
let swap (a: array int) (i: int) (j: int) =
    let v = a[i] in
    a[i] <- a[j];
    a[j] <- v</pre>
```

```
let selection_sort (a: array int) =
  for i = 0 to length a - 1 do
    let imin = ref i in
    for j = i + 1 to length a - 1 do
        invariant { i <= !imin < j }
        invariant { forall k: int. i <= k < j -> a[!imin] <= a[k] }
        if a[j] < a[!imin] then imin := j
        done;</pre>
```

swap a !min i

done



Why3 (5/8)

• theories on arrays

```
let swap (a: array int) (i: int) (j: int) =
   requires { 0 <= i < length a /\ 0 <= j < length a }
   ensures { exchange (old a) a i j }
   let v = a[i] in
   a[i] <- a[j];
   a[j] <- v</pre>
```

(see the why3 libraries)

http://why3.lri.fr

Why₃ (6/8)

theories on arrays

```
let selection_sort (a: array int) =
    ensures { sorted a \land permut (old a) a }
'L:
    for i = 0 to length a - 1 do
      invariant { sorted_sub a 0 i \land permut (at a 'L) a}
      invariant { forall k1 k2: int. 0 <= k1 < i <= k2 < length a -> a[k1] <= a[k2] }
      let imin = ref i in
      for j = i + 1 to length a - 1 do
        invariant { i <= !imin < j }</pre>
        invariant { forall k: int. i <= k < j -> a[!imin] <= a[k] }
        if a[j] < a[!imin] then imin := j
      done;
      swap a !imin i ;
    done
                                      imin
          0
```



Why3 (7/8)

- interfaces with automatic provers (SMT's)
- SMT tool successful if «good assertion»
 - impact on writings of Hoare logic formulae
 - impact on program text
- Alt-Ergo among best [LRI, Conchon, et al]
- Z3 is excellent [MSRR, Bjorner/de Moura]
- CVC3 top on recursive datatypes
- Gappa for real numbers [Inria, Melquiond]

Why3 (8/8)

- interfaces with interactive proof assistants
- PVS [SRI, Shankar]
- Coq [Inria, Herbelin et al]
 - Why3 theories are translated to Coq
 - lengthy proofs are feasible
 - use SSreflect commands to shorten proofs [MSR-Inria, Gonthier
 et al]
 - unfortunately Why3 is not fully compatible with SSreflect

Demos

A few sorting algorithms

- demos
- insertion sort

A few sorting algorithms

• quicksort

Conclusions

Conclusion (1/3)

- Automatic part of proof for tedious case analyzes
- Interactive proofs for the conceptual part of the algorithm

- From interactive part, one must call the automatic part
 - possible extensions of Why3 theories
 - but typing problems (inside Coq)

Conclusion (2/3)

- Hoare logic prevents to write awkward denotational semantics
- Nobody cares about termination !

- Explore simple programs about algorithms before jumping to large programs.
- Why3 memory model is naive. It is a «back-end for other systems».
- Plan to experiment on graph algorithms and prove all Sedgewick's book on algorithms.

Conclusion (3/3)

- Why3 is **excellent** for mixing formal proofs and SMT's calls
- Interface still rough for beginners
- Concurrency ?
- Functional programs ?
- Hoare logic vs Type refinements (F* [MSR])
- Frama-C project at french CEA extends Why3 to C programs.