

| | | | |
|---------------|-------|-------------------------|--|
| var, x | | term variable | |
| loc, l | | location | |
| $tvar, X$ | | type variable | |
| $term, t$ | $::=$ | | term |
| | | x | variable |
| | | $\lambda x . t$ | bind x in t abstraction |
| | | $t_1 t_2$ | application |
| | | $()$ | unit |
| | | l | location |
| | | ref t | reference |
| | | $!t$ | get |
| | | $t_1 := t_2$ | set |
| | | (t) | S |
| | | $\{ t_1 / x \} t_2$ | M |
| $value, v$ | $::=$ | | value |
| | | $\lambda x . t$ | bind x in t abstraction |
| | | $()$ | unit |
| | | l | location |
| $type, S, T$ | $::=$ | | type |
| | | X | variable |
| | | $T \rightarrow T'$ | function |
| | | unit | unit |
| | | T ref | ref |
| | | (T) | S |
| env, Γ | $::=$ | | type environment |
| | | empty | |
| | | $env, x : T$ | |
| phi, Φ | $::=$ | | type location environment |
| | | empty | |
| | | $phi, l : T$ | |
| sto, μ | $::=$ | | store |
| | | empty | |
| | | $sto, l : t$ | |
| $formula$ | $::=$ | | |
| | | $judgement$ | |
| | | $x : T$ in env | |
| | | $l : T$ in phi | |
| | | $l : t$ in sto | |
| | | ok env | |
| | | $l \# \mu$ | |
| $terminals$ | $::=$ | | |
| | | λ | |
| | | \rightarrow | |

| | | | |
|--------------------|-----|-------------------|--|
| | | \vdash | |
| | | \longrightarrow | |
| <i>Stook</i> | ::= | | ok sto |
| <i>Jtype</i> | ::= | | $\Gamma! \Phi \vdash t : T$ |
| <i>Jop</i> | ::= | | $t_1 \mu_1 \longrightarrow t_2 \mu_2$ $t_1 \mu_1$ reduces to $t_2 \mu_2$ |
| <i>judgement</i> | ::= | | <i>Stook</i> |
| | | | <i>Jtype</i> |
| | | | <i>Jop</i> |
| <i>user_syntax</i> | ::= | | <i>var</i> |
| | | | loc |
| | | | <i>tvar</i> |
| | | | <i>term</i> |
| | | | <i>value</i> |
| | | | <i>type</i> |
| | | | <i>env</i> |
| | | | <i>phi</i> |
| | | | <i>sto</i> |
| | | | <i>formula</i> |
| | | | <i>terminals</i> |

ok sto

$\frac{}{\mathbf{ok\ empty}}$ STO_OK_EMPTY

$\frac{\mathbf{ok\ } \mu}{\mathbf{ok\ } \mu, l : t}$ STO_OK_PUSH

$\Gamma! \Phi \vdash t : T$

$\frac{\mathbf{ok\ } \Gamma \quad x : T \mathbf{in\ } \Gamma}{\Gamma! \Phi \vdash x : T}$ TYPING_VAR

$\frac{\Gamma, x : S! \Phi \vdash t : T}{\Gamma! \Phi \vdash \lambda x. t : S \rightarrow T}$ TYPING_ABS

$\frac{\Gamma! \Phi \vdash t_1 : S \rightarrow T \quad \Gamma! \Phi \vdash t_2 : S}{\Gamma! \Phi \vdash t_1 t_2 : T}$ TYPING_APP

$\frac{\mathbf{ok\ } \Gamma}{\Gamma! \Phi \vdash () : \mathbf{unit}}$ TYPING_UNIT

$$\begin{array}{c}
\text{ok}\Gamma \\
\hline
\Gamma! \Phi \vdash l : T \mathbf{ref} \quad \text{TYPING_LOC} \\
\hline
\Gamma! \Phi \vdash t_1 : T \\
\hline
\Gamma! \Phi \vdash \mathbf{ref} t_1 : T \mathbf{ref} \quad \text{TYPING_NEW} \\
\hline
\Gamma! \Phi \vdash t_1 : T \mathbf{ref} \\
\hline
\Gamma! \Phi \vdash !t_1 : T \quad \text{TYPING_GET} \\
\hline
\Gamma! \Phi \vdash t_1 : T \mathbf{ref} \\
\Gamma! \Phi \vdash t_2 : T \\
\hline
\Gamma! \Phi \vdash t_1 := t_2 : \mathbf{unit} \quad \text{TYPING_SET}
\end{array}$$

$$\boxed{t_1 | \mu_1 \longrightarrow t_2 | \mu_2}$$

$t_1 | \mu_1$ reduces to $t_2 | \mu_2$

$$\frac{\text{ok}\mu}{(\lambda x . t_1) v_2 | \mu \longrightarrow \{v_2 / x\} t_1 | \mu} \quad \text{RED_BETA}$$

$$\frac{\text{ok}\mu \quad l \# \mu}{\mathbf{ref} v | \mu \longrightarrow l | \mu, l : v} \quad \text{RED_NEW}$$

$$\frac{\text{ok}\mu \quad l : t \mathbf{in} \mu}{!l | \mu \longrightarrow t | \mu} \quad \text{RED_GET}$$

$$\frac{\text{ok}\mu}{l := v | \mu \longrightarrow () | \mu, l : v} \quad \text{RED_SET}$$

$$\frac{t_1 | \mu \longrightarrow t'_1 | \mu'}{t_1 t_2 | \mu \longrightarrow t'_1 t_2 | \mu'} \quad \text{RED_APP_1}$$

$$\frac{t_2 | \mu \longrightarrow t'_2 | \mu'}{v_1 t_2 | \mu \longrightarrow v_1 t'_2 | \mu'} \quad \text{RED_APP_2}$$

$$\frac{t_1 | \mu \longrightarrow t'_1 | \mu'}{\mathbf{ref} t_1 | \mu \longrightarrow \mathbf{ref} t'_1 | \mu'} \quad \text{RED_NEW_1}$$

$$\frac{t_1 | \mu \longrightarrow t'_1 | \mu'}{!t_1 | \mu \longrightarrow !t'_1 | \mu'} \quad \text{RED_GET_1}$$

$$\frac{t_1 | \mu \longrightarrow t'_1 | \mu'}{t_1 := t_2 | \mu \longrightarrow t'_1 := t_2 | \mu'} \quad \text{RED_SET_1}$$

$$\frac{t_2 | \mu \longrightarrow t'_2 | \mu'}{v_1 := t_2 | \mu \longrightarrow v_1 := t'_2 | \mu'} \quad \text{RED_SET_2}$$

Definition rules: 20 good 0 bad
Definition rule clauses: 45 good 0 bad